# RGFuzz: Rule-Guided Fuzzer for WebAssembly Runtimes

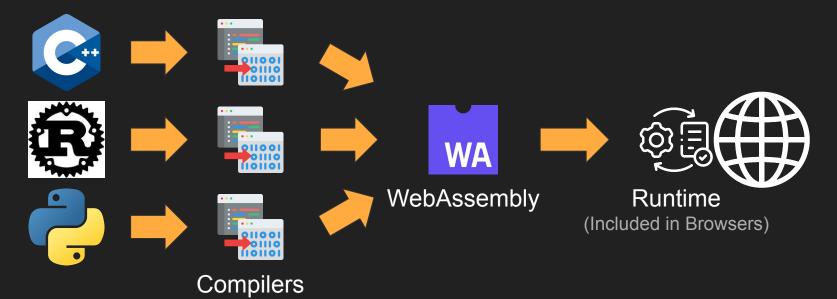
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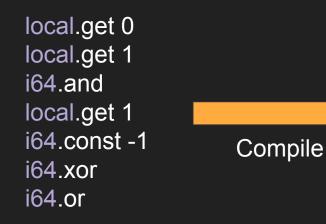
## WebAssembly (WASM)

- Fast, safe, portable, and compact language
- Best for compilation target for other languages



#### WebAssembly Runtimes

- WebAssembly runs on a stack machine
- Stack machine is slow  $\rightarrow$  Let's compile the code!
- Just-In-Time (JIT) compilation to machine code



push rbp mov rbp, rsp not rcx mov rax, rdx or rax, rcx mov rsp, rbp pop rbp ret

- Optimizations to further boost speed

local.get 0 local.get 1 i64.and local.get 1 i64.const -1 i64.xor i64.or



1. Translate to IR

[Args] v0: i64, v1: i64

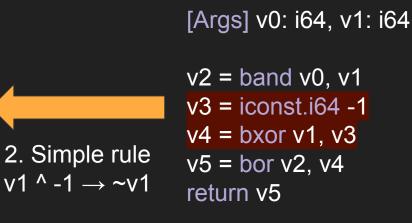
- Apply simple rule: v1 ^ -1  $\rightarrow$  ~v1 (changed to not)

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[Args] v0: i64, v1: i64

v2 = band v0, v1

v4 = bnot v1 v5 = bor v2, v4 return v5 2. Simple rule  $v1^{-1} \rightarrow \sim v1$ 

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- Can also apply complex rule: (v0 & v1) |  $\sim$ v1  $\rightarrow$  v0 |  $\sim$ v1

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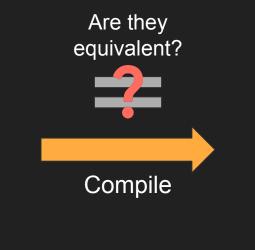
[Args] v0: i64, v1: i64 v2 = band v0, v1 v4 = bnot v1 v5 = bor v2, v4 return v5 ( $v0 \otimes v1$ ) | ~v1  $\rightarrow v0$  | ~v1 (Args] v0: i64, v1: i64 v2 = v0 v4 = bnot v1 v5 = bor v2, v4 return v5

## Semantic Bugs

- What happens if optimization rules are wrongly written?

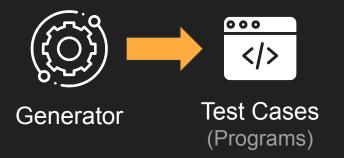
- Semantic bug: For some input, exec. of original code != exec. of compiled code

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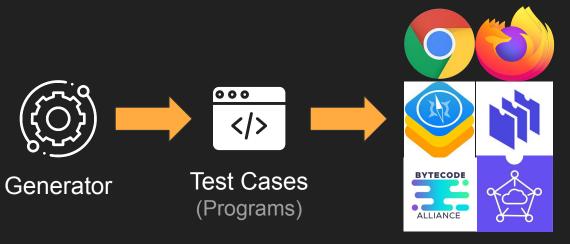


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- Differential fuzzing



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#### Execution

- Differential fuzzing



test cases efficiently?

- Differential fuzzing



- Challenge 1.1: Complex rules

 $(v0 \& v1) | \sim v1$  $\rightarrow v0 | \sim v1$ 

Testing needs:

Optimization

local.get 0 local.get 1 i64.and local.get 1 i64.const -1 i64.xor i64.or

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[Preliminary Study] SOTA fuzzers failed to generate such program (Xsmith, wasm-smith)

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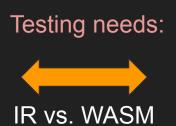
#### [Challenge]

Compiler rules: Defined in IR Programs: Written in WASM

- Challenge 1.1: Complex rules

 $(v0 \& v1) | \sim v1$  $\rightarrow v0 | \sim v1$ 

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#### [Challenge]

Compiler rules: Defined in IR Programs: Written in WASM

How do we close the gap??

## Approach 1.1: Instruction-level Inference

- Challenge 1.2: Closing the gap between IR and WebAssembly

Instruction-level Inference:

band  $\rightarrow$  i64.and bor  $\rightarrow$  i64.or bnot  $\rightarrow$  ???

### Approach 1.1: Instruction-level Inference

- Challenge 1.2: Closing the gap between IR and WebAssembly

Instruction-level Inference:

band  $\rightarrow$  i64.and bor  $\rightarrow$  i64.or bnot  $\rightarrow$  ??? We do not have a WASM instruction that directly maps to bnot

### Approach 1.2: Rule-level Inference

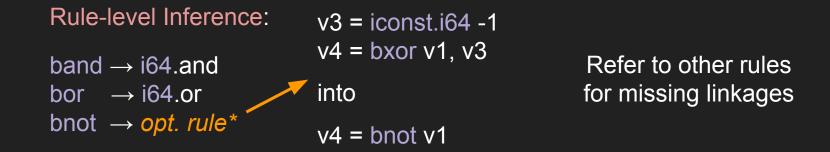
- Challenge 1.2: Closing the gap between IR and WebAssembly

Rule-level Inference:

band  $\rightarrow$  i64.and bor  $\rightarrow$  i64.or bnot  $\rightarrow$  opt. rule\* Refer to other rules for missing linkages

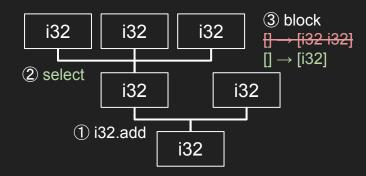
### Approach 1.2: Rule-level Inference

- Challenge 1.2: Closing the gap between IR and WebAssembly



- Challenge 2: Generate structures or instructions diversely

- AST-based : limited structure diversity (e.g., blocks)
- Stack-based: limited instruction diversity (e.g., select)



**AST-based** 

Instructions	Stack
① local.get 0	0
2 local.get 1	[i32]
③ i32.add <del>select</del>	[i32 i32]
(d) block	[i32]
[] → [i32 i32]	[i32 i32 i32]

Stack-based

#### - Solution: Reverse stack-based generation

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e.g., v128.bitselect requires 3 v128s on parameters, but only 1 in returns

## Evaluation

- Target Runtimes: 6 runtimes
  - wasmtime, Wasmer, WasmEdge, V8, SpiderMonkey, JavaScriptCore
  - Tested various optimization / architectures

- Found 20 new bugs, with one CVE ID (CVE-2023-29548)

## Evaluation

- Coverage

- Able to cover significantly more in wasmtime

Baseline	Coverage	RGFUZZ Mean	Mean	RGFUZZ <sup>-</sup> Difference	p	wasr Mean	ntime-differen Difference	tial p	Mean	Fuzzgen Difference	p	Mean	Wasm-mutate Difference	p	Mean	Xsmith Difference	p
Optimization	Line	69.37%	49.29%	20.08%	0.008	36.89%	32.48%	0.008	46.32%	23.05%	0.008	37.85%	31.52%	0.008	39.33%	30.04%	0.008
	Rule	70.64%	51.89%	18.75%	0.012	42.79%	27.85%	0.012	50.97%	19.67%	0.008	40.13%	30.51%	0.008	41.02%	29.62%	0.012
Lowering	Line	71.43%	70.54%	0.89%	0.008	65.42%	6.01%	0.008	56.61%	14.82%	0.008	27.22%	44.21%	0.008	27.53%	43.90%	0.008
	Rule	75.11%	74.23%	0.88%	0.012	68.73%	6.38%	0.008	56.63%	18.48%	0.008	27.06%	48.05%	0.008	26.47%	48.64%	0.012
Total	Line	28.14%	27.12%	1.02%	0.008	17.41%	10.73%	0.012	16.46%	11.68%	0.008	11.85%	16.29%	0.012	20.37%	7.77%	0.008

## Evaluation

#### - Coverage

- Able to cover significantly more in wasmtime
- Could also efficiently test other runtimes
- Rule-guided fuzzing was only effective in wasmtime though

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Engines	RGFuzz Mean	Mean	RGFUZZ <sup>-</sup> Difference	p	Mean	Wasm-smith Difference	p	Mean	Xsmith Difference	p
Wasmer WasmEdge V8	<b>12.22%</b> 10.82% <b>19.81%</b>	12.21% <b>10.83%</b> 19.80%	0.01% -0.01% 0.01%	0.234 0.154 0.264	12.02% 10.57% 19.53%	0.19% 0.25% 0.28%	0.012 0.011 0.011	11.33% 10.37% 17.65%	0.89% 0.45% 2.16%	0.011 0.010 0.011
JavaScriptCore	19.58%	19.55%	0.03%	0.057	19.07%	0.51%	0.012	14.58%	5.00%	0.012

## Case Study

#### - Could even cover super complex optimizations

```
(rule (simplify (bor ty @ $I64
(bor ty
    (bor ty
    (ishl ty x (iconst_u ty 56))
    (ishl ty
         (band ty x (iconst_u ty 0xff00))
        (iconst_u ty 40)))
    (bor ty
    (ishl ty
         (band ty x (iconst_u ty 0xff_0000))
         (iconst_u ty 24))
    (ishl ty
         (band ty x (iconst_u ty 0xff00_0000))
         (iconst_u ty 8))))
(bor ty
    (bor ty
    (band ty
         (ushr ty x (iconst_u ty 8))
        (iconst_u ty 0xff00_0000))
    (band ty
         (ushr ty x (iconst_u ty 24))
        (iconst_u ty 0xff_0000)))
    (bor ty
    (band ty
         (ushr ty x (iconst_u ty 40))
         (iconst_u ty 0xff00))
    (ushr ty x (iconst_u ty 56)))))
(bswap tv x))
```

**Complex Optimization Rule** 

**Specific Immediates** 

## Case Study

- Could find optimization bugs!

local.get	0		
local.get	1		
local.get	1		
local.get	0		
i64.gt_s			
select			
;; Expecte	ed:	<pre>min(arg0,</pre>	arg1)
;; Actual	:	<pre>max(arg0,</pre>	arg1)

min mistaken as max wasmtime issue 8114

100	al.get	2	;;	arg2:	0	xffi	E8			
f64	.load									
f64	.const	0								
100	al.get	0	;;	arg0:	: 1					
sel	ect									
;;	Expecte	ed:	m	em[arg	]2]	or	0.0	based	on	argo
;;	Actual		T	rap						0.0000000000000000000000000000000000000

Load more bytes than expected (as xmm) wasmtime issue 8112

## Key Takeaways

- Two main approaches
  - Rule-guided fuzzing
  - Reverse stack-based generation
- Showed effectiveness in finding optimization bugs